1.1 Data structures

Data structures

A ***data structure*** is a way of organizing, storing, and performing operations on data. Operations performed on a data structure include accessing or updating stored data, searching for specific data, inserting new data, and removing data. The following provides a list of basic data structures.

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Table 1.1.1: Basic data structures.

Data structure	Description
Record	A ***record*** is the data structure that stores subitems, often called fields, with a name associated with each subitem.
Array	An ***array*** is a data structure that stores an ordered list of items, where each item is directly accessible by a positional index.
Linked list	A ***linked list*** is a data structure that stores an ordered list of items in nodes, where each node stores data and has a pointer to the next node.
Binary tree	A ***binary tree*** is a data structure in which each node stores data and has up to two children, known as a left child and a right child.
Hash table	A ***hash table*** is a data structure that stores unordered items by mapping (or hashing) each item to a location in an array.
Heap	A ***max-heap*** is a tree that maintains the simple property that a node's key is greater than or equal to the node's childrens' keys. A ***min-heap*** is a tree that maintains the simple property that a node's key is less than or equal to the node's childrens' keys.
Graph	A ***graph*** is a data structure for representing connections among items, and consists of vertices connected by edges. A ***vertex*** represents an item in a graph. An ***edge*** represents a connection between two vertices in a graph.

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PARTICIPATION ACTIVITY 1.1.1: Basic data structures.

1) A linked list stores items in an unspecified order.
   - True
   - False

2) A node in binary tree can have zero, one, or two children.
   - True
   - False

3) A list node's data can store a record with multiple subitems.
   - True
   - False

4) Items stored in an array can be accessed using a positional index.
   - True
   - False

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Choosing data structures

The selection of data structures used in a program depends on both the type of data being stored and the operations the program may need to perform on that data. Choosing the best data structure often requires determining which data structure provides a good balance given expected uses. Ex: If a program requires fast insertion of new data, a linked list may be a better choice than an array.

PARTICIPATION ACTIVITY 1.1.2: A linked list avoids the shifting problem.

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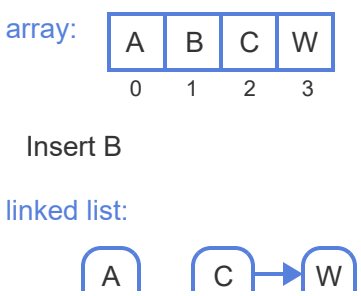

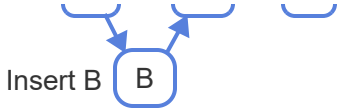

Animation content:

Static figure: Pictured are an array and a linked list. Each data structure is accompanied by the text Insert B.

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The array has these elements in this order: A, B, C, W. Each of these four elements is numbered according to its position in the array; that is, A is numbered 0, B is numbered 1, C is numbered 2, and W is numbered 3.

The linked list has the same four elements in the same order: A, B, C, W. The node containing element A has an arrow pointing down-right to the node with element B, which has an arrow pointing up-right to the node with C, which has an arrow pointing right to the last node with W.

Step 1: Inserting an item at a specific location in an array requires making room for the item by shifting higher-indexed items.
Array with contents [A, C, W] appears, with the elements numbered 0, 1, 2 respectively. The text Insert B also appears. In the array, an empty slot appears to the right of the slot containing element W and is numbered 3. W and C each shift 1 position to the right to make room for a new array entry at index 1.

Step 2: Once the higher index items have been shifted, the new item can be inserted at the desired index.
The B from the text Insert B replicates and the replica moves to the now empty slot in the array at index 1. The array thus contains these elements: A, B, C, W

Step 3: To insert a new item in a linked list, a list node for the new item is first created.
A linked list appears with nodes, each node having one element. The node containing element A points to the node containing element C, which points to the node containing element W. A new node appears that contains element B.

Step 4: Item B's next pointer is assigned to point to item C. Item A's next pointer is updated to point to item B. No shifting of other items was required.
An arrow appears that points from the node with element B to the node with C; then, the arrow pointing from the node with element A to the node with B moves such that it now points to the node with element B.

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Animation captions:

1. Inserting an item at a specific location in an array requires making room for the item by shifting higher-indexed items.
2. Once the higher index items have been shifted, the new item can be inserted at the desired index.
3. To insert new item in a linked list, a list node for the new item is first created.
4. Item B's next pointer is assigned to point to item C. Item A's next pointer is updated to point to item B. No shifting of other items was required.

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PARTICIPATION ACTIVITY 1.1.3: Basic data structures.

1) Inserting an item at the end of a 999-item array requires how many items to be shifted?

Check Show answer

2) Inserting an item at the end of a 999-item linked list requires how many items to be shifted?

Check Show answer

3) Inserting an item at the beginning of a 999-item array requires how many items to be shifted?

Check Show answer

4) Inserting an item at the beginning of a 999-item linked list requires how many items to be shifted?

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Check Show answer

1.2 Introduction to algorithms

Algorithms

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An ***algorithm*** describes a sequence of steps to solve a computational problem or perform a calculation. An algorithm can be described in English, pseudocode, a programming language, hardware, etc. A ***computational problem*** specifies an input, a question about the input that can be answered using a computer, and the desired output.

PARTICIPATION ACTIVITY 1.2.1: Computational problems and algorithms.

Computational problem:

Input: Array of integers

13	7	45	-3	92	17

Question: What is the maximum value in the input array?

Output: Maximum integer in input array
92

Algorithm:

```
FindMax(inputArray) {
   max = inputArray[0]

   for (i = 1; i < inputArray.size; ++i) {
      if (inputArray[i] > max) {
         max = inputArray[i];
      }
   }

   return max
}
```

Animation content:

Static figure: Pictured are a computational problem and a corresponding algorithm to solve the problem.

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Computational problem text:
Input: Array of integers: 13, 7, 45, -3, 92, 17
Question: What is the maximum value in the input array?
Output: Maximum integer in input array 92
End computational problem text.

Algorithm:
Begin Pseudocode:

```
FindMax(inputArray) {
  max = inputArray[0]

  for (i = 1; i < inputArray.size; ++i) {
    if (inputArray[i] > max) {
      max = inputArray[i;]
    }
  }
  return max
}
End Pseudocode.
```

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Step 1: A computational problem is a problem that can be solved using a computer. A computational problem specifies the problem input, a question to be answered, and the desired output.
A skeleton for a computation problem appears. It has three labeled sections, Input, Question, and Output. These sections are empty.

Step 2: For the problem of finding the maximum value in an array, the input is an array of numbers. In the Input section of the computational problem, the text Array of integers appears alongside an array of integers: 13, 7, 45, -3, 92, 17.

Step 3: The question problem is: What is the maximum value in the input array? The problems output is a single value that is the maximum value in the array.
The output also includes the answer, 92.

Step 4: The FindMax algorithm defines a sequence of steps that determines the maximum value in the array.
The code from the static figure now appears.

Animation captions:

1. A computational problem is a problem that can be solved using a computer. A computational problem specifies the problem input, a question to be answered, and the desired output.
2. For the problem of finding the maximum value in an array, the input is an array of numbers.
3. The problem's question is: What is the maximum value in the input array? The problem's output is a single value that is the maximum value in the array.
4. The FindMax algorithm defines a sequence of steps that determines the maximum value in the array.

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PARTICIPATION ACTIVITY 1.2.2: Algorithms and computational problems.

Consider the problem of determining the number of times (or frequency) a specific word appears in a list of words.

1) Which can be used as the problem input?

- String for user-specified word
- Array of unique words and string for user-specified word
- Array of all words in the list and string for user-specified word

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2) What is the problem output?

- Integer value for the frequency of most frequent word
- String value for the most frequent word in input array
- Integer value for the frequency of specified word

3) An algorithm to solve this computation problem must be written using a programming language.

- True
- False

Practical applications of algorithms

Computational problems can be found in numerous domains, including e-commerce, internet technologies, biology, manufacturing, transportation, etc. Algorithms have been developed for numerous computational problems within these domains.

A computational problem can be solved in many ways, but finding the best algorithm to solve a problem can be challenging. However, many computational problems have common subproblems, for which efficient algorithms have been developed. The examples below describe a computational problem within a specific domain and list a common algorithm (each discussed elsewhere) that can be used to solve the problem.

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Table 1.2.1: Example computational problems and common algorithms.

Application domain	Computational problem	Common algorithm
DNA analysis	Given two DNA sequences from different individuals, what is the longest shared sequence of nucleotides?	*Longest common substring problem:* A longest common substring algorithm determines the longest common substring that exists in two input strings. DNA sequences can be represented using strings consisting of the letters A, C, G, and T to represent the four different nucleotides.
Search engines	Given a product ID and a sorted array of all in-stock products, is the product in stock and what is the product's price?	*Binary search:* The binary search algorithm is an efficient algorithm for searching a list. The list's elements must be sorted and directly accessible (such as an array).
Navigation	Given a user's current location and desired location, what is the fastest route to walk to the destination?	*Dijkstra's shortest path:* Dijkstra's shortest path algorithm determines the shortest path from a start vertex to each vertex in a graph. The possible routes between two locations can be represented using a graph, where vertices represent specific locations and connecting edges specify the time required to walk between those two locations.

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PARTICIPATION ACTIVITY 1.2.3: Computational problems and common algorithms.

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Match the common algorithm to another computational problem that can be solved using that algorithm.

If unable to drag and drop, refresh the page.

Binary search **Shortest path algorithm** **Longest common substring**

	Do two student essays share a common phrase consisting of a sequence of more than 100 letters?
	Given the airports at which an airline operates and distances between those airports, what is the shortest total flight distance between two airports?
	Given a sorted list of a company's employee records and an employee's first and last name, what is a specific employee's phone number?

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Reset

Efficient algorithms and hard problems

Computer scientists and programmers typically focus on using and designing efficient algorithms to solve problems. Algorithm efficiency is most commonly measured by the algorithm runtime, and an efficient algorithm is one whose runtime increases no more than polynomially with respect to the input size. However, some problems exist for which an efficient algorithm is unknown.

NP-complete problems are a set of problems for which no known efficient algorithm exists. NP-complete problems have the following characteristics:

- No efficient algorithm has been found to solve an NP-complete problem.
- No one has proven that an efficient algorithm to solve an NP-complete problem is impossible.
- If an efficient algorithm exists for one NP-complete problem, then all NP-complete problems can be solved efficiently.

By knowing a problem is NP-complete, instead of trying to find an efficient algorithm to solve the problem, a programmer can focus on finding an algorithm to efficiently find a good, but non-optimal, solution.

PARTICIPATION ACTIVITY 1.2.4: Example NP-complete problem: Cliques.

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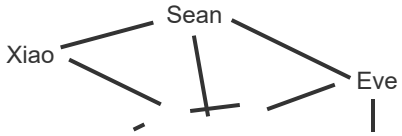
Computational problem:

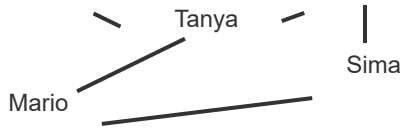
Input:
G: Graph of social network, where vertices represent individuals and edges represent mutual friendship
K: integer

Input social network:



Question: Does a set of K people who all know each other exist?

Output:
Boolean value



Problem is NP-complete

Does a set of 3 people who all know each other exist? Yes
Does a set of 4 people who all know each other exist? No

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Animation content:

Static figure:
Pictured are a computational problem and an input social network.

Computational problem:
Input: "G: Graph of social network, where vertices represent individuals and edges represent mutual friendship Boolean value Problem is NP-complete. K: integer"
Question: Does a set of K people who all know each other exist?
Output: Boolean value

This problem is NP-complete

Six names exist in the graph, some of which are connected to each other with an edge. The following table gives a source-vertex name, followed by a colon, with all the other names connected to the source-vertex name after the colon. (e.g., Xiao is connected to Sean and Tanya)
Xiao: Sean, Tanya
Sean: Xiao, Tanya, Eve
Eve: Sean, Tanya, Sima
Tanya: Xiao, Sean, Eve, Mario
Mario: Tanya, Sima
Sima: Eve, Mario

Two labels exist in the lower-right:
"Does a set of 3 people who all know each other exist? Yes"
"Does a set of 4 people who all know each other exist? No"

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Step 1: A programmer may be asked to write an algorithm to solve the problem of determining if a set of K people who all know each other exists within a graph of a social network.
The computational problem visible in the static figure appears.

Step 2: For the example social network graph and K = 3, the algorithm should return yes. Xiao, Sean, and Tanya all know each other. Sean, Tanya, and Eve also all know each other.
The input social network graph appears along with the question, "Does a set of 3 people who all

know each other exist?" The names in the graph Xiao, Sean, Tanya, are highlighted, making note that the three all know each other. The question now is answered with the text "Yes."

Step 3: For K = 4, no set of 4 individuals who all know each other exists, and the algorithm should return no.
Some question and answer text appears: "Does a set of 4 people who all know each other exist? No"

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Step 4: This problem is equivalent to the clique decision problem, which is NP-complete, and no known polynomial time algorithm exists.

Animation captions:

1. A programmer may be asked to write an algorithm to solve the problem of determining if a set of K people who all know each other exists within a graph of a social network.
2. For the example social network graph and K = 3, the algorithm should return yes. Xiao, Sean, and Tanya all know each other. Sean, Tanya, and Eve also all know each other.
3. For K = 4, no set of 4 individuals who all know each other exists, and the algorithm should return no.
4. This problem is equivalent to the clique decision problem, which is NP-complete, and no known polynomial time algorithm exists.

PARTICIPATION ACTIVITY 1.2.5: Efficient algorithm and hard problems.

1) An algorithm with a polynomial runtime is considered efficient.
   - ○ True
   - ○ False

2) An efficient algorithm exists for all computational problems.
   - ○ True
   - ○ False

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3) An efficient algorithm to solve an NP-complete problem may exist.
   - ○ True
   - ○ False

1.3 Relation between data structures and algorithms

Algorithms for data structures

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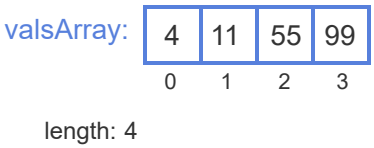
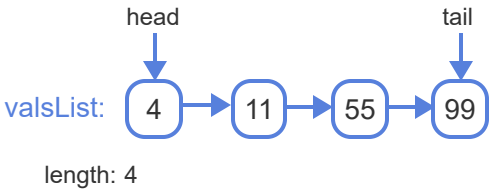
Data structures not only define how data is organized and stored, but also the operations performed on the data structure. While common operations include inserting, removing, and searching for data, the algorithms to implement those operations are typically specific to each data structure. Ex: Appending an item to a linked list requires a different algorithm than appending an item to an array.

PARTICIPATION ACTIVITY 1.3.1: A list avoids the shifting problem.

Append(valsArray, 99)



Append(valsList, 99)



Algorithm for appending to array

```
ArrayAppend(array, newItem) {
   currentSize = array→length
   Increase array size by one
   array[currentSize] = newItem
}
```

Algorithm for appending to linked-list

```
ListAppend(list, newNode) {
   if (list→head == null) { // List empty
      list→head = newNode
      list→tail = newNode
   }
   else{
      list→tail→next = newNode
      list→tail = newNode
   }
   list→length++
}
```

Animation content:

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Static figure: An array and a linked list, each shown with an algorithm for appending an item.

On the left is a code statement Append(valsArray, 99). An array named valsArray is shown, with contents: 4, 11, 55, 99. Text notes that the length is 4.
Below the array is a code block labeled Algorithm for appending to array.
Begin pseudocode:
ArrayAppend(array, newItem) {

```
    currentSize = array···→length
    Increase array size by one
    array[currentSize] = newItem
}
End pseudocode.
```

On the right is the code statement Append(valsList, 99). A linked list named valsList is shown, with 4 nodes. Head points to a node with value 4, which points to a node with value 11, which points to a node with value 55, which points to a node with value 99. Tail points to the node with value 99. Text notes that the length is 4.

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Below the linked list is a code block labeled Algorithm for appending to linked list.

```
Begin pseudocode:
ListAppend(list, newNode) {
  if (list···→head == null) { // List empty
    list···→head = newNode
    list···→tail = newNode
  }
  else{
    list···→tail···→next = newNode
    list···→tail = newNode
  }
  list···→length++
}
End pseudocode.
```

Step 1: The algorithm to append an item to an array determines the current size, increases the array size by 1, and assigns the new item as the last array element. valsArray starts with length 3 and values 4, 11, and 55 at indices 0 to 2 respectively. A box is added to the array at index 3. Length changes to 4, and the value 99 appears in the new box.

Step 2: The algorithm to append an item to a linked list points the next pointer of the tail node and the tail pointer of the list to the new node. The linked list starts with length 3. Head points to value 4, which points to value 11, which points to value 55. Tail points to 55. A new node with value 99 appears next to the node with value 55. A pointer is added from 55 to 99, and tail moves to point to 99. Length is updated to 4.

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Animation captions:

1. The algorithm to append an item to an array determines the current size, increases the array size by 1, and assigns the new item as the last array element.
2. The algorithm to append an item to a linked list points the tail node's next pointer and the list's tail pointer to the new node.

PARTICIPATION ACTIVITY 1.3.2: Algorithms for data structures.

Consider the array and linked list in the animation above. Can the following algorithms be implemented with the same code for both an array and linked list?

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1) Append an item
   - Yes
   - No

2) Return the first item
   - Yes
   - No

3) Return the current size
   - Yes
   - No

Algorithms using data structures

Some algorithms utilize data structures to store and organize data during the algorithm execution. Ex: An algorithm that determines a list of the top five salespersons, may use an array to store salespersons sorted by their total sales.

Figure 1.3.1: Algorithm to determine the top five salespersons using an array.

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```
DisplayTopFiveSalespersons(allSalespersons) {
   // topSales array has 5 elements
   // Array elements have subitems for name and total sales
   // Array will be sorted from highest total sales to lowest total sales
   topSales = Create array with 5 elements

   // Initialize all array elements with a negative sales total
   for (i = 0; i < topSales⇢length; ++i) {
      topSales[i]⇢name = ""
      topSales[i]⇢salesTotal = -1
   }

   for each salesPerson in allSalespersons {
      // If salesPerson's total sales is greater than the last
      // topSales element, salesPerson is one of the top five so far
      if (salesPerson⇢salesTotal > topSales[topSales⇢length -
1]⇢salesTotal) {

         // Assign the last element in topSales with the current
salesperson
         topSales[topSales⇢length - 1]⇢name =  salesPerson⇢name
         topSales[topSales⇢length - 1]⇢salesTotal =
salesPerson⇢salesTotal

         // Sort topSales in descending order
         SortDescending(topSales)
      }
   }

   // Display the top five salespersons
   for (i = 0; i < topSales⇢length; ++i) {
      Display topSales[i]
   }
}
```

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PARTICIPATION ACTIVITY 1.3.3: Top five salespersons.

1) Which of the following is *not* equal to the number of items in the topSales array?

   - topSales⇢length
   - 5
   - allSalesperson⇢length

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2) To adapt the algorithm to display the top 10 salespersons, what modifications are required?

- Only the array creation
- All loops in the algorithm
- Both the array creation and all loops

3) If allSalespersons has length 3 then DisplayTopFiveSalespersons() displays five elements, two of which have no name and -1 sales.
   - True
   - False

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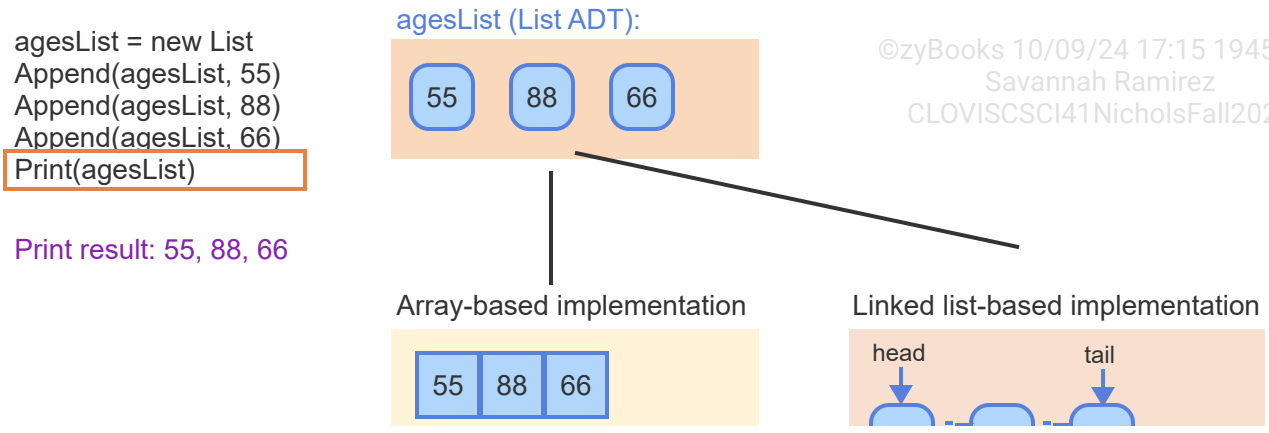
1.4 Abstract data types

Abstract data types (ADTs)

An ***abstract data type*** (***ADT***) is a data type described by predefined user operations, such as "insert data at rear," without indicating how each operation is implemented. An ADT can be implemented using different underlying data structures. However, a programmer need not have knowledge of the underlying implementation to use an ADT.

Ex: A list is a common ADT for holding ordered data, having operations like append a data item, remove a data item, search whether a data item exists, and print the list. A list ADT is commonly implemented using arrays or linked list data structures.

PARTICIPATION ACTIVITY 1.4.1: List ADT using array and linked lists data structures.



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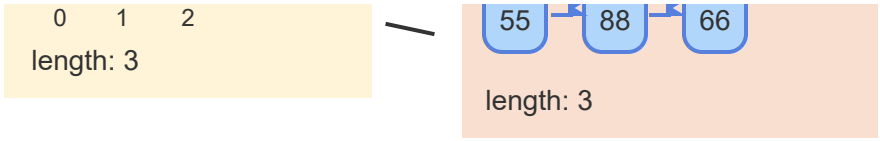

Animation content:

Static figure: Code to create a list and append values 55, 88, and 66 to a list is shown in upper left.
Begin Pseudocode:
agesList = new List
Append(agesList, 55)
Append(agesList, 88)
Append(agesList, 66)
Print(agesList)
End Pseudocode.

The last line of the code, Print(agesList), is boxed in. Under the code, there is some text: Print result: 55, 88, 66

agesList shown in upper center as an abstract data type (List ADT). The ADT is represented as three items with values 55, 88, and 66 respectively. Below agesList ADT, two possible implementations are shown.
First implementation uses an array with items: 55, 88, and 66, and a length variable equal to value 3.
Second implementation uses a linked list with nodes 55 pointing to 88 pointing to 66, and a length variable equal to 3. 55 is in the head node and 66 is in the tail node.

Step 1: A new list named agesList is created. Items can be appended to the list. The items are ordered. The first four lines of code, which initialize agesList and append to it three times, result in a List with three elements that is shown as an ADT, without implementation details.

Step 2: Printing the list prints the items in order. The last line of the code, Print(agesList) is highlighted and the text, Print result: 55, 88, 66, appears.

Step 3: A list ADT is commonly implemented using array and linked list data structures. But, a programmer need not have knowledge of which data structure is used to use the list ADT. Two lines extrude from the List ADT representation; one line connects the representation to the Array-based implementation and the other to the Linked-list-based implementation.

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Animation captions:

1. A new list named agesList is created. Items can be appended to the list. The items are ordered.
2. Printing the list prints the items in order.

3. A list ADT is commonly implemented using array and linked list data structures. But, a programmer need not have knowledge of which data structure is used to use the list ADT.

PARTICIPATION ACTIVITY 1.4.2: Abstract data types.

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1) Starting with an empty list, what is the list contents after the following operations?

   Append(list, 11)
   Append(list, 4)
   Append(list, 7)

   - 4, 7, 11
   - 7, 4, 11
   - 11, 4, 7

2) A remove operation for a list ADT removes the specified item. Given a list with contents: 2, 20, 30, what is the list after the following operation?

   Remove(list, item 2)

   - 2, 30
   - 2, 20, 30
   - 20, 30

3) A programmer must know the underlying implementation of the list ADT in order to use a list.

   - True
   - False

4) A list ADT's underlying data structure has no impact on the program's execution.

   - True
   - False

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Common ADTs

Table 1.4.1: Common ADTs.

Abstract data type	Description	Common underlying data structures
List	A ***list*** is an ADT for holding ordered data.	Array, linked list
Dynamic array	A ***dynamic array*** is an ADT for holding ordered data and allowing indexed access.	Array
Stack	A ***stack*** is an ADT in which items are only inserted on or removed from the top of a stack.	Linked list
Queue	A ***queue*** is an ADT in which items are inserted at the end of the queue and removed from the front of the queue.	Linked list
Deque	A ***deque*** (pronounced "deck" and short for double-ended queue) is an ADT in which items can be inserted and removed at both the front and back.	Linked list
Bag	A ***bag*** is an ADT for storing items in which the order does not matter and duplicate items are allowed.	Array, linked list
Set	A ***set*** is an ADT for a collection of distinct items.	Binary search tree, hash table
Priority queue	A ***priority queue*** is a queue where each item has a priority, and items with higher priority are closer to the front of the queue than items with lower priority.	Heap
Dictionary (Map)	A ***dictionary*** is an ADT that associates (or maps) keys with values.	Hash table, binary search tree

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PARTICIPATION ACTIVITY 1.4.3: Common ADTs.

Consider the ADTs listed in the table above. Match the ADT with the description of the order and uniqueness of items in the ADT.

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Set | Bag | List | Priority queue

	Items are ordered based on how items are added. Duplicate items are allowed.
	Items are not ordered. Duplicate items are not allowed.
	Items are ordered based on items' priority. Duplicate items are allowed.
	Items are not ordered. Duplicate items are allowed.

Reset

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1.5 Applications of ADTs

Abstraction and optimization

Abstraction means to have a user interact with an item at a high-level, with lower-level internal details hidden from the user. ADTs support abstraction by hiding the underlying implementation details and providing a well-defined set of operations for using the ADT.

Using abstract data types enables programmers or algorithm designers to focus on higher-level operations and algorithms, thus improving programmer efficiency. However, knowledge of the underlying implementation is needed to analyze or improve the runtime efficiency.

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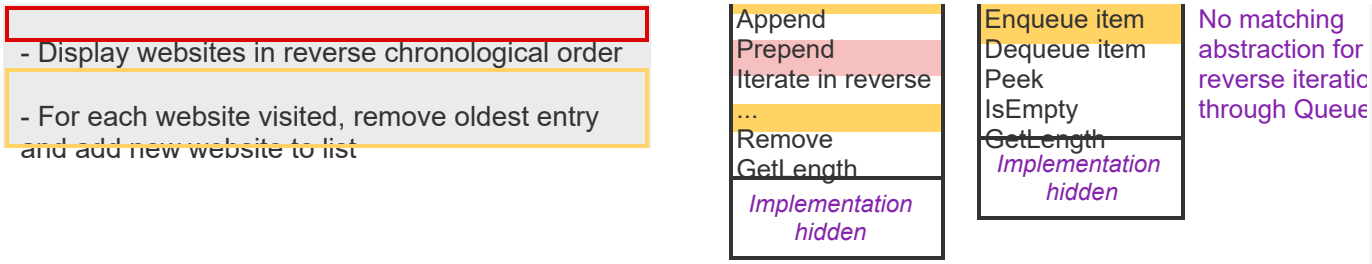
PARTICIPATION ACTIVITY 1.5.1: Programming using ADTs.

Program requirements

- Maintain list of 5 most recently visited websites

Available ADTs

List | Queue



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Animation content:

Static figure:
Program requirements are listed on the left. Two are highlighted:
- "Maintain list of 5 most recently visited websites"
- "Display websites in reverse chronological order", which is highlighted red
- "For each website visited, remove oldest entry and add new website to list", which is highlighted yellow

There are two available ADTs, List and Queue, on the right. Available operations for each ADT are below that ADT's name. Some operations are colored in corresponding ways to the highlights on the requirements.
List has these operations: Append, which is yellow; Prepend; Iterate in reverse, which is red; an ellipsis here designates that there are more operations now shown; Remove, which is yellow; GetLength. Below the last operation is the text "Implementation hidden."

Queue has these operations: Enqueue item, which is yellow; Dequeue item, which is yellow; Peek; IsEmpty; GetLength. Below the last operation is the text, "Implementation hidden." There is also text that reads, "No matching abstraction for reverse iteration through Queue."

Step 1: Abstraction simplifies programming. ADTs allow programmers to focus on choosing which ADTs best match a program's needs. The requirements and the two available ADTs are visible.

Step 2: Both the list and queue ADTs support efficient interfaces for removing items from one end (removing oldest entry) and adding items to the other end (adding new entries). The requirement, "For each website visited, remove oldest entry and add new website to list," is highlighted yellow. In the List ADT, the "Append" and "Remove" operations are colored yellow. In the Queue ADT, the "Enqueue item" and "Dequeue item" operations are colored yellow.

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Step 3: The list ADT supports iterating through list contents in reverse order, but the queue ADT does not. The requirement, "Display websites in reverse chronological order," is highlighted red. In the List ADT, the "Iterate in reverse" operation is colored red. Next to the Queue ADT, the text, "No matching abstraction for reverse iteration through Queue," appears.

Step 4: To use the list (or queue) ADT, the programmer does not need to know the underlying

implementation. Below each set of operations for both of the ADTs is the text, "Implementation hidden."

Animation captions:

1. Abstraction simplifies programming. ADTs allow programmers to focus on choosing which ADTs best match a program's needs.
2. Both the list and queue ADTs support efficient interfaces for removing items from one end (removing oldest entry) and adding items to the other end (adding new entries).
3. The list ADT supports iterating through list contents in reverse order, but the queue ADT does not.
4. To use the list (or queue) ADT, the programmer does not need to know the underlying implementation.

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PARTICIPATION ACTIVITY 1.5.2: Programming with ADTs.

Consider the example in the animation above.

1) The _____ ADT is the better match for the program's requirements.
   - queue
   - list

2) The list ADT _____.
   - can only be implemented using an array
   - can only be implemented using a linked list
   - can be implemented in numerous ways

3) Knowledge of an ADT's underlying implementation is needed to analyze the runtime efficiency.
   - True
   - False

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ADTs in standard libraries

Most programming languages provide standard libraries that implement common abstract data types. Some languages allow programmers to choose the underlying data structure used for the ADTs. Other programming languages may use a specific data structure to implement each ADT, or may automatically choose the underlying data-structure.

Table 1.5.1: Standard libraries in various programming languages.

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Programming language	Library	Common supported ADTs
Python	Python standard library	list, set, dict, deque
C++	Standard template library (STL)	vector, list, deque, queue, stack, set, map
Java	Java collections framework (JCF)	Collection, Set, List, Map, Queue, Deque

PARTICIPATION ACTIVITY 1.5.3: ADTs in standard libraries.

1) Python, C++, and Java all provide built-in support for a deque ADT.
   - True
   - False

2) The underlying data structure for a list ADT is the same for all programming languages.
   - True
   - False

3) ADTs are only supported in standard libraries.
   - True
   - False

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1.6 Algorithm efficiency

Algorithm efficiency

An algorithm describes the method to solve a computational problem. Programmers and computer scientists should use or write efficient algorithms. ***Algorithm efficiency*** is typically measured by the algorithm's computational complexity. ***Computational complexity*** is the amount of resources used by the algorithm. The most common resources considered are the runtime and memory usage.

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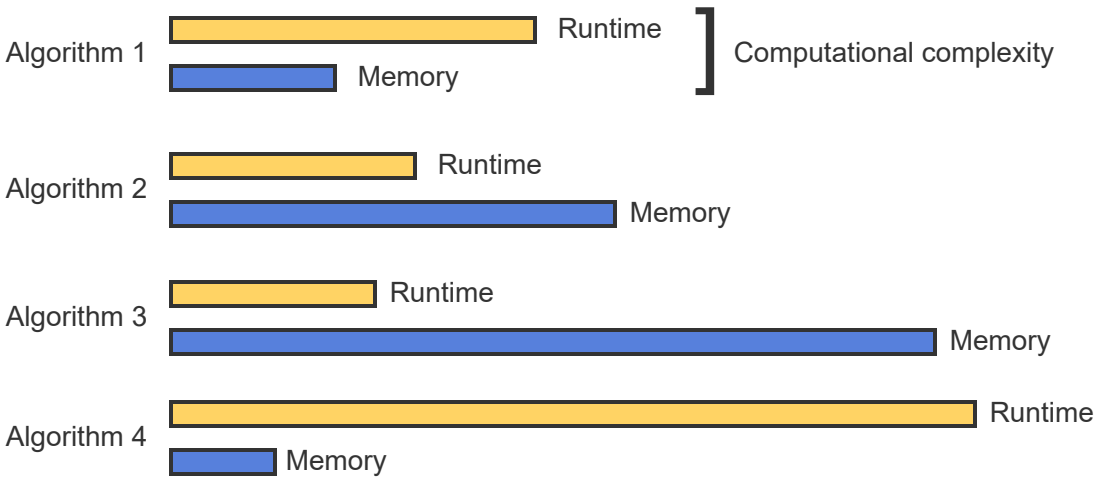
PARTICIPATION ACTIVITY 1.6.1: Computational complexity.



Animation content:

Static figure:
Runtime is shown for four algorithms. Each has a bar for the runtime and memory, which together comprise the algorithm's computational complexity. Longer bars indicate longer runtime or higher memory usage.
Runtime bars ordered from shortest to longest: algorithm 3, algorithm 2, algorithm 1, algorithm 4.
Memory bars order from shortest to longest: algorithm 4, algorithm 1, algorithm 2, algorithm 3.

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Animation captions:

1. An algorithm's computational complexity includes runtime and memory usage.
2. Measuring runtime and memory usage allows different algorithms to be compared.
3. Complexity analysis is used to identify and avoid using algorithms with long runtimes or high memory usage.

PARTICIPATION ACTIVITY 1.6.2: Algorithm efficiency and computational complexity.

1) Computational complexity analysis allows the efficiency of algorithms to be compared.
   - True
   - False

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2) Two different algorithms that produce the same result have the same computational complexity.
   - True
   - False

3) Runtime and memory usage are the only two resources making up computational complexity.
   - True
   - False

Runtime complexity, best case, and worst case

An algorithm's ***runtime complexity*** is a function, T(N), that represents the number of constant time operations performed by the algorithm on an input of size N. Runtime complexity is discussed in more detail elsewhere.

Because an algorithm's runtime may vary significantly based on the input data, a common approach is to identify best and worst case scenarios. An algorithm's ***best case*** is the scenario where the algorithm does the minimum possible number of operations. An algorithm's ***worst case*** is the scenario where the algorithm does the maximum possible number of operations.

Input data size must remain a variable

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A best case or worst case scenario describes contents of the algorithm's input data only. The input data size must remain a variable, N. Otherwise, the overwhelming majority of algorithms would have a best case of N=0, since no input data would be processed. In both theory and practice, saying "the best case is when the algorithm

doesn't process any data" is not useful. Complexity analysis always treats the input data size as a variable.

PARTICIPATION ACTIVITY 1.6.3: Linear search best and worst cases.

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```
LinearSearch(numbers, numbersSize, key) {
  i = 0
  while (i < numbersSize) {
     if (numbers[i] == key)
        return i
     i = i + 1
  }

  return -1 // not found
}
```

numbers:

54	79	26	91	29	33

key = 26: neither best nor worst case

key = 54: best case

key = 82: worst case

Animation content:

Static figure: The array contains [54, 79, 26, 91, 29, 33]. LinearSearch returns the index where the given key is found or -1 if the key is not found.
The search for key 26 is neither best nor worst case.
The search for key 54 is best case.
The search for key 82 is worst case.
Begin Pseudocode:

```
LinearSearch(numbers, numbersSize, key) {
   i = 0
   while (i < numbersSize) {
      if (numbers[i] == key)
         return i
      i = i + 1
   }

   return -1 // not found
}
```

End Pseudocode.

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Step 1: LinearSearch searches through array elements until finding the key. Searching for 26 requires iterating through the first 3 elements.

Step 2: The search for 26 is neither the best nor the worst case.

Step 3: Searching for 54 only requires one comparison and is the best case: The key is found at the start of the array. No other search could perform fewer operations.

Step 4: Searching for 82 compares against all array items and is the worst case: The number is not found in the array. No other search could perform more operations.

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Animation captions:

1. LinearSearch searches through array elements until finding the key. Searching for 26 requires iterating through the first 3 elements.
2. The search for 26 is neither the best nor the worst case.
3. Searching for 54 only requires one comparison and is the best case: The key is found at the start of the array. No other search could perform fewer operations.
4. Searching for 82 compares against all array items and is the worst case: The number is not found in the array. No other search could perform more operations.

PARTICIPATION ACTIVITY 1.6.4: FindFirstLessThan algorithm best and worst case.

Consider the following function that returns the first value in a list that is less than the specified value. If no list items are less than the specified value, the specified value is returned.

```
FindFirstLessThan(list, listSize, value) {
   for (i = 0; i < listSize; i++) {
      if (list[i] < value)
         return list[i]
   }
   return value // no lesser value found
}
```

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Neither best nor worst case **Worst case** **Best case**

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	No items in the list are less than `value`.
	The first half of the list has elements greater than `value` and the second half has elements less than `value`.

The first item in the list is less than `value`.

Reset

PARTICIPATION ACTIVITY 1.6.5: Best and worst case concepts.

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1) The linear search algorithm's best case scenario is when N = 0.
   - True
   - False

2) An algorithm's best and worst case scenarios are always different.
   - True
   - False

Space complexity

An algorithm's ***space complexity*** is a function, S(N), that represents the number of fixed-size memory units used by the algorithm for an input of size N. Ex: The space complexity of an algorithm that duplicates a list of numbers is S(N) = 2N + k, where k is a constant representing memory used for things like the loop counter and list pointers.

Space complexity includes the input data and additional memory allocated by the algorithm. An algorithm's ***auxiliary space complexity*** is the space complexity not including the input data. Ex: An algorithm to find the maximum number in a list will have a space complexity of S(N) = N + k, but an auxiliary space complexity of S(N) = k, where k is a constant.

PARTICIPATION ACTIVITY 1.6.6: FindMax space complexity and auxiliary space complexity.

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```
FindMax(list, listSize) {
   if (listSize >= 1) {
      maximum = list[0]
      i = 1
      while (i < listSize) {
         if (list[i] > maximum) {
            maximum = list[i]
         }
         i = i + 1
      }
      return maximum
```

Memory (input)

listSize: fixed-size integer

list: Array of N integers

Memory (non-input)

maximum: fixed-size integer

```
    }
}
```

i: fixed-size integer

2 fixed-size variables

Space complexity: S(N) = N + 3

Auxiliary space complexity: S(N) = 2

Animation content:

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Static figure:

FindMax implementation is shown with an array argument (list) and an integer argument (listSize), and two local variables (maximum and i). Memory is categorized into two parts: input and non-input. Input memory includes list and listSize. Non-input memory includes variables i and maximum.

Begin Pseudocode:

```
FindMax(list, listSize) {
   if (listSize >= 1) {
      maximum = list[0]
      i = 1
      while (i < listSize) {
         if (list[i] > maximum) {
            maximum = list[i]
         }
         i = i + 1
      }
      return maximum
   }
}
```

End Pseudocode.

Step 1: FindMax arguments represent input data. Non-input data includes variables allocated in the function body: maximum and i.

Step 2: The list size is a variable, N. Three integers are also used, listSize, maximum, and i, making the space complexity S(N) = N + 3.

Step 3: The auxiliary space complexity includes only the non-input data, which does not increase for larger input lists.

Step 4: The function auxiliary space complexity is S(N) = 2.

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Animation captions:

1. FindMax's arguments represent input data. Non-input data includes variables allocated in the function body: maximum and i.
2. The list's size is a variable, N. Three integers are also used, listSize, maximum, and i, making the space complexity S(N) = N + 3.

3. The auxiliary space complexity includes only the non-input data, which does not increase for larger input lists.
4. The function's auxiliary space complexity is S(N) = 2.

PARTICIPATION ACTIVITY 1.6.7: Space complexity of GetEvens function.

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Consider the following function, which builds and returns a list of even numbers from the input list.

```
GetEvens(list, listSize) {
   i = 0
   evensList = Create new, empty list
   while (i < listSize) {
      if (list[i] % 2 == 0)
         Add list[i] to evensList
      i = i + 1
   }
   return evensList
}
```

1) What is the maximum possible size of the returned list?
   - ○ listSize
   - ○ listSize / 2

2) What is the minimum possible size of the returned list?
   - ○ listSize / 2
   - ○ 1
   - ○ 0

3) What is the worst case auxiliary space complexity of GetEvens if N is the list's size and k is a constant?
   - ○ S(N) = N + k
   - ○ S(N) = k

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4) What is the best case auxiliary space complexity of GetEvens if N is the list's size and k is a constant?
   - ○ S(N) = N + k
   - ○ S(N) = k